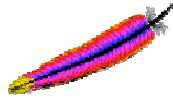


Advanced Topics in Module Design: Threading and Portability

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I'm a member of the Apache Software Foundation, and a contributor to the HTTPD and APR projects. I've also been known to dabble in other open source projects, such as OpenSLP and PHP.



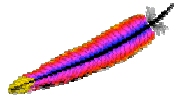
Thread-safe

From The Free On-line Dictionary of Computing (09 FEB 02) [foldoc]:

thread-safe

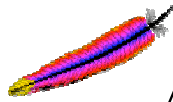
A description of code which is either *re-entrant* or protected from multiple simultaneous execution by some form of *mutual exclusion*.

(1997-01-30)



The Basics

Some APR Primitive Types



A Who's Who of Mutexes

- `apr_thread_mutex_t`
- `apr_proc_mutex_t`
- `apr_global_mutex_t`

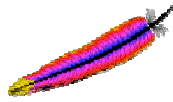


- `apr_foo_mutex_lock()`
 - Grab the lock, or block until available
- `apr_foo_mutex_unlock()`
 - Release the current lock

The `apr_thread_mutex_t` type is used to synchronize execution between threads in the same process. The `apr_proc_mutex_t` type is used to synchronize different processes, but not necessarily multiple threads within those same processes. Finally, the `apr_global_mutex_t` type ensures that multiple threads in multiple processes are all properly synchronized.

One typically uses `apr_thread_mutex_t` to protect access to a resource in a single process, while `apr_global_mutex_t` can be used for resources that are shared between multiple multithreaded processes (which is typical for an MPM-agnostic module). One case where an `apr_proc_mutex_t` may be optimal would be when each process can guarantee that only one thread at a time will access some central resource. An example of this can be seen with the `socket accept()` code inside Apache 2.0.

Other than creating and destroying mutexes, the only operations that can be performed on them are the `_lock()` and `_unlock()` calls.



Reader/Writer Locks

- **`apr_thread_rwlock_t`**

- `apr_thread_rwlock_rdlock()`

- Grab the shared read lock, blocks for any writers

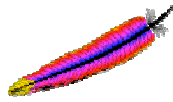
- `apr_thread_rwlock_wrlock()`

- Grab the exclusive write lock, blocking new readers

- `apr_thread_rwlock_unlock()`

- Release the current lock

Reader/Writer locks are created and destroyed in much the same way mutexes are, but behave quite differently. Instead of having a simple locked/unlocked state, a reader/writer lock allows the caller to signify whether they desire a shared read state, or an exclusive write state of the lock. This allows finer-grain control over the synchronization between threads, which can lead to better parallelism. The way this works is you can have multiple readers at the same time, but as soon as a writer requests the lock, no new readers are allowed to acquire until that writer has released, and that writer does not acquire the lock until all the current readers have released.



Condition Variables

- **apr_thread_cond_t**

- `apr_thread_cond_wait()`

- Sleep until any signal arrives

- `apr_thread_cond_signal()`

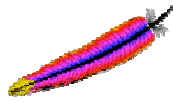
- Send a signal to one waiting thread

- `apr_thread_cond_broadcast()`

- Send a signal to all waiting threads

Condition variables are always used in conjunction with a mutex. This mutex protects the signaling mechanism from race condition problems.

One common example that illustrates a use of condition variables is the producer/consumer problem. Imagine a single producer that can produce 5 widgets per second. There are also 10 consumers, each of which can consume 1 widget per second. Given only mutexes, there is no guarantee that the work will be shared between all of the consumers. With condition variables, each consumer waits for the signal to begin work. As the producer finishes readying each widget, he signals one consumer, who then wakes up and performs a consumptive task. Upon completing the consumption, the worker then returns to the queue.



Threads



- **apr_thread_t**

- `apr_thread_create()`

- Create a new thread (with specialized attributes)

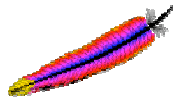
- `apr_thread_exit()`

- Exit from the current thread (with a return value)

- `apr_thread_join()`

- Block until a particular thread calls
`apr_thread_exit()`

Thread creation is fairly straightforward. Each thread is created from a *pool*, from which a thread-specific sub-pool is created and passed to the start function. A start function is specified at creation time, and as soon as the thread is created by the system, it begins executing that start function. APR provides a data passing system from the creating thread to the created thread, via the start function's prototype.



One-time Calls

- **apr_thread_once_t**

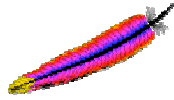
- `apr_thread_once_init()`

- Initialize an `apr_thread_once_t` variable

- `apr_thread_once()`

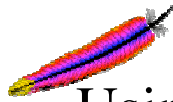
- Execute the given function once

One-time calls allow initialization of process-wide resources in a way that avoids race conditions. The call to `apr_thread_once()` guarantees that a specified call will only be called once, no matter how many times `apr_thread_once()` is called.



Apache 2.0 Hooks

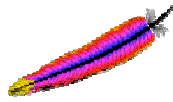
Using mutexes and shared memory
within the Apache framework



Using APR Primitives in Apache

- mutexes
- reader/writer locks
- condition variables
- shared memory
- ...

The way we create, initialize, use, and destroy each of these types is pretty much the same across the board. This means that you can write the code once and expect it to work on both Unix and Windows, BeOS and Netware. APR tries to hide much of the trouble of dealing with inconsistencies between the platforms.

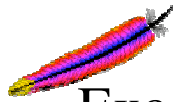


Global Mutex Creation

- Create it in the Parent:
 - Usually in `post_config` hook
- Attach to it in the Child:
 - This is the `child_init` hook

By creating in the parent process, we allow systems like Unix to inherit these types in a child process.

The `post_config` hook is ideal In most cases for module authors, since it is the first hook in the parent where all the configuration information is available. For example, a module may define a directive where a user can specify a filename where a global mutex is to be stored.



Example: Create a Global Mutex

```
static int shm_counter_post_config(apr_pool_t *pconf,
                                   apr_pool_t *plog,
                                   apr_pool_t *ptemp,
                                   server_rec *s) {

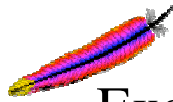
    int rv;
    shm_counter_scfg_t *scfg;

    /* Get the module configuration */
    scfg = ap_get_module_config(s->module_config,
                                &shm_counter_module);

    /* Create a global mutex from the config directive */
    rv = apr_global_mutex_create(&scfg->mutex,
                                scfg->shmcounterlockfile,
                                APR_LOCK_DEFAULT, pconf);
}
```

This example shows how we can pull the module configuration from the `server_rec` and use those directives as parameters to the global mutex creation routine.

Passing `APR_LOCK_DEFAULT` here signifies that we wish to have a non-recursive mutex. A recursive mutex is a mutex that can be repeatedly acquired by the same thread without immediately deadlocking (as long as the same number of releases occur). I do not believe that there are many cases that require a recursive lock, and tend to avoid them. In addition, they tend to perform poorly on some Unix systems, and also tend to mask programmer bugs more often than non-recursive mutexes.



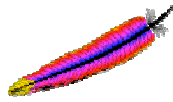
Example: Attach Global Mutex

```
static void shm_counter_child_init(apr_pool_t *p,  
                                   server_rec *s)  
{  
    apr_status_t rv;  
    shm_counter_scfg_t *scfg = ap_get_module_config(  
        s->module_config,  
        &shm_counter_module);  
  
    /* Now that we are in a child process, we have to  
     * reconnect to the global mutex. */  
    rv = apr_global_mutex_child_init(&scfg->mutex,  
                                     scfg->shmcounterlockfile,  
                                     p);  
}
```

In the `child_init` hook we must attach to the mutex. Under Unix, this happens automatically, since mutexes are inherited by child processes via the `fork()` mechanism.

The reason we need to pass the filename of the mutex back into the `apr_global_mutex_child_init()` call is so that under Windows we are able to connect to the same global mutex from all child processes.

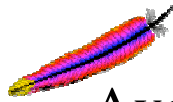
At this point, any child process may lock or unload the global mutex in order to synchronize access to some cross-process shared resource.



Common Pitfall

- The double DSO-load problem
 - Apache loads each module twice:
 - First time to see if it fails at startup
 - Second time to actually load it
 - (They are loaded again upon each restart...)

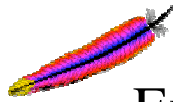
Apache wants to be sure that all modules can be successfully loaded into the server before the server decides it is ok to start up. To do this it goes through a full load-unload-load cycle at startup. The problem with this is that it causes unexpected problems when we try to initialize our mutex twice.



Avoiding the Double DSO-load

- Solution:
 - Don't create mutexes during the first load
- 1. First time in `post_config` we set a flag
- 2. Next time through we look for that flag
 - if it is set, we create the mutex

One way to prevent this from happening is to set a flag on the first call to the `post_config` hook. On each subsequent call to `post_config` we check if that flag has been set, and if so we proceed with our initialization.



Example: Double DSO-load

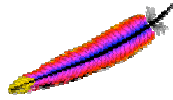
```
static int shm_counter_post_config(apr_pool_t *pconf, apr_pool_t *plog,
                                   apr_pool_t *ptemp, server_rec *s)
{
    apr_status_t rv;
    void *data = NULL;
    const char *userdata_key = "shm_counter_post_config";

    apr_pool_userdata_get(&data, userdata_key, s->process->pool);
    if (data == NULL) {
        /* WARNING: This must *not* be apr_pool_userdata_setn(). */
        apr_pool_userdata_set((const void *)1, userdata_key,
                              apr_pool_cleanup_null, s->process->pool);
        return OK; /* This would be the first time through */
    }

    /* Proceed with normal mutex and shared memory creation . . . */
}
```

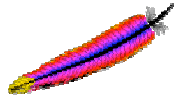
Bonus Question: Why is it so important that we don't use `apr_pool_userdata_setn()` to set the userdata flag?

The reason for this is because the static symbol section of the DSO may not be at the same address offset when it is reloaded. Since `setn()` does not make a copy and only compares addresses, the `get()` will be unable to find the original userdata.



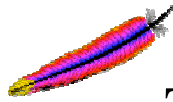
Summary

1. Create in the Parent (`post_config`)
 2. Attach in the Child (`child_init`)
- This works for these types:
 - mutexes
 - condition variables
 - reader/writer locks
 - shared memory
 - etc...



Shared Memory

Efficient and portable shared memory
for your module

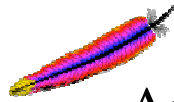


Types of Shared Memory

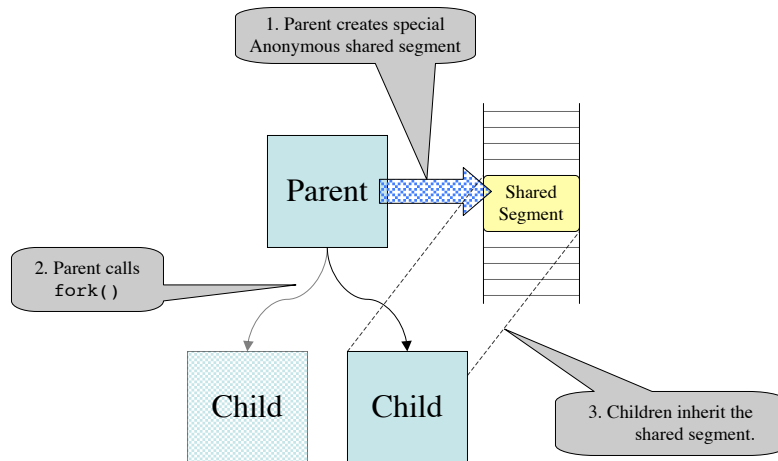
- Anonymous
 - Requires process inheritance
 - Created in the parent
 - Automatically inherited in the child
- Name-based
 - Associated with a file
 - Processes need not be ancestors
 - Must deal with file permissions

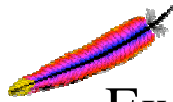
The main difference between the two has to do with process inheritance. Anonymous shared memory is typically easier to deal with, but is not as well supported as name-based shared memory. Windows does not support anonymous shared memory (with some exceptions).

Often a particular implementation of APR's anonymous shared memory will have size limitations, while that same platform may not have similar restrictions on a name-based implementation.



Anonymous Shared Memory





Example: Anonymous Shmem

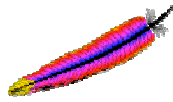
```
static int shm_counter_post_config(apr_pool_t *pconf,
                                   apr_pool_t *plog,
                                   apr_pool_t *ptemp,
                                   server_rec *s)
{
    int rv;
    ...

    /* Create an anonymous shared memory segment by passing
     * a NULL as the shared memory filename */
    rv = apr_shm_create(&scfg->counters_shm,
                       sizeof(*scfg->counters),
                       NULL, pconf);
}
```

As with the mutex example, this example shows how we can pull the module configuration from the `server_rec` and use that information to create our segment.

Note that the double DSO-load preventing code is omitted for brevity, as is the mutex creation code.

By passing the requested size to the creation function we are guaranteed by APR that it will create at least that size segment or the create routine will fail. There is a small amount of overhead in some cases, but it is handled in a transparent manner.

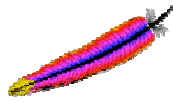


Accessing the Segment

- Segment is mapped as soon as it is created
 - It has a start address
 - You can query that start address
- Reminder: The segment may not be mapped at the same address within different processes.

```
scfg->counters = apr_shm_baseaddr_get(scfg->counters_shm);
```

Accessing the segment is as simple as asking APR for the address of the start of the segment. Note that this address is only valid for the current process, not across processes. What this means is that simple pointers can not be stored in the segment, but must be relative to some global addressing scheme. A pointer in one process will not necessarily work in another process where the segment may have been mapped at a totally different location.



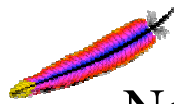
Windows Portability

- Windows can't inherit shared memory
 - it has no `fork()` call!
- Stupid Windows
- Solution:
 - Just like with mutexes:
 - Attach from the “child” process
(hint: to be portable to Windows, we can only use name-based shared memory.)

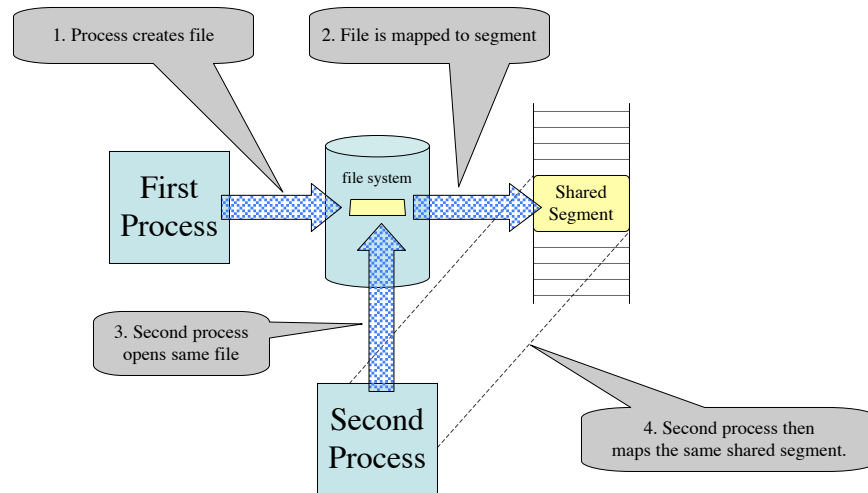
This isn't entirely true, there are ways around this, and there are types of “anonymous shared memory” available on Windows, but in the end there still needs to be a way for the Apache “parent” and “child” to rendezvous so that they can both attach to the same memory segment.

The best way to avoid all of this is to just use name-based shared memory. Unfortunately this means that you have to start worrying about things like where to put the file, and what permissions can it have.

The Apache 2.0 scoreboard is an excellent example of a hybrid solution. If the underlying system support anonymous shared memory, the scoreboard uses it, otherwise it falls back to name-based shared memory, using the name of a file from a configuration directive, or defaulting to a standard filename.

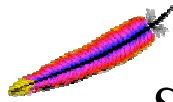


Name-based Shared Memory



Bonus Question: Why do we have to worry about race conditions?

The second process cannot access the segment until after the first process has completely created it.



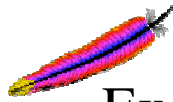
Sharing with external apps

- Must use name-based
- Associate it with a file
- Other programs attach to that file

- Beware of race conditions
- Beware of permission problems
(note recent security problem in Apache scoreboard)

Although APR's name-based shared memory requires that the segment be associated with a file, whether or not it uses that file for the actual contents of the segment is implementation (and platform) specific. By associating the segment with a file, APR allows fully autonomous external programs to all share the same segment. By this property of name-based shared memory, it is possible to write an external application that can monitor the status of your module via its file-based shared memory segment.

Recently there was a security problem in Apache (both 1.3 and 2.0) where it was possible for someone who had access to the user under which apache was running to send certain signals to any process on the system as though the signals were coming from root. The reason for this was that the apache user was able to open the shared memory segment and alter the entries holding the PIDs of the children, causing the parent process (running as root) to unwittingly send signals to whatever PID it read from the segment. Ouch!



Example: Name-based Shmem

```
static int shm_counter_post_config(apr_pool_t *pconf,
                                   apr_pool_t *plog,
                                   apr_pool_t *ptemp,
                                   server_rec *s) {

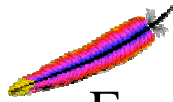
    int rv;
    shm_counter_scfg_t *scfg;
    ...

    /* Get the module configuration */
    scfg = ap_get_module_config(s->module_config,
                                &shm_counter_module);

    /* Create a name-based shared memory segment using the filename
     * out of our config directive */
    rv = apr_shm_create(&scfg->counters_shm, sizeof(*scfg->counters),
                        scfg->shmcounterfile, pconf);
}
```

Note again that the double DSO-load preventing code is omitted for brevity, as is the mutex creation code.

The only difference here from how we created the anonymous shared memory segment is that now we are passing an actual filename instead of just NULL.



Example: Name-based Shmem (cont)

```
static void shm_counter_child_init(apr_pool_t *p,  
                                   server_rec *s)  
{  
    apr_status_t rv;  
    shm_counter_scfg_t *scfg  
        = ap_get_module_config(s->module_config,  
                               &shm_counter_module);  
  
    rv = apr_shm_attach(&scfg->counters_shm,  
                       scfg->shmcounterfile, p);  
  
    scfg->counters = apr_shm_baseaddr_get(scfg->counters_shm);  
}
```

Now that we have one process which has created and (automatically) attached to the segment, we can have other programs use that same filename to attach to the same segment.

Once we have attached to the segment, we can grab that base address to the segment and start plugging away.

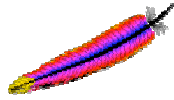


RMM (Relocatable Memory Manager)

- Provides `malloc()` and `free()`
- Works with any block of memory
- Estimates overhead
- Threadsafe
- Usable on shared memory segments

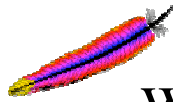
The RMM was created specifically for memory management within any contiguous block of memory, and is particularly suited for memory management within shared memory space. It is threadsafe in the sense that it can be accessed in parallel by multiple threads and processes without memory corruption.

There is a significant amount of overhead associated with maintain a free list within the shared-memory segment, but that cost can be estimated by the RMM code itself.



Apache 2.0 Architecture

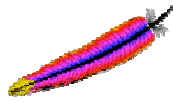
A quick MPM overview



What's new in Apache 2.0?

- Filters
- MPMs
- SSL
- lots more...

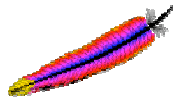
The main thing that we are concerned with in terms of writing threadsafe and multithread-aware modules is the MPMs.



What is an MPM?

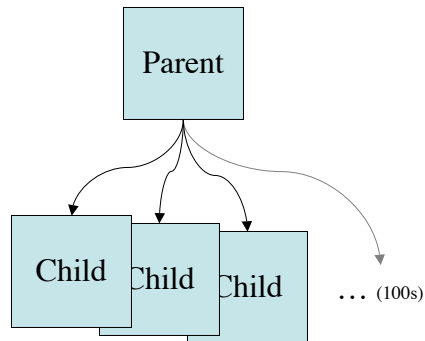
- “Multi-processing Module”
 - Different HTTP server process models
 - Each give us
 - Platform-specific features
 - Admin may chose suitable:
 - Reliability
 - Performance
 - Features

An MPM allows platform-specific optimizations in how Apache deals with incoming requests and parallelism.

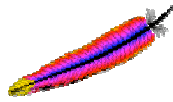


Prefork MPM

- Classic Apache 1.3 model
- 1 connection per Child
- Pros:
 - Isolates faults
 - Performs well
- Cons:
 - Scales poorly (high memory reqts.)

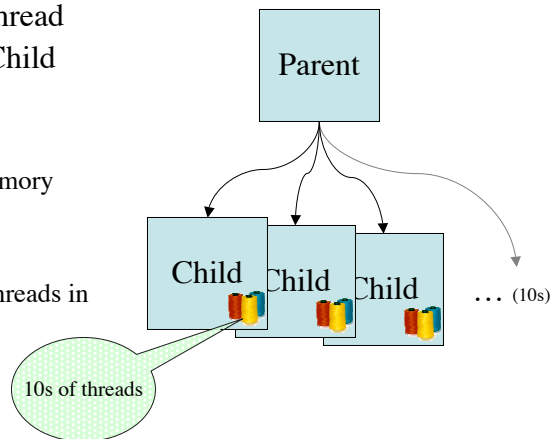


Available on Unix. This MPM matches the model used since NCSA and through Apache up to Apache 1.3. In general it tends to perform well on uniprocessor machines, and is highly robust. The prefork MPM is not well suited for high-traffic, highly concurrent systems. Although the prefork MPM processes many hundreds of simultaneous connections without problem, systems are usually limited by the amount of memory consumed.

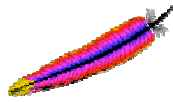


Worker MPM

- Hybrid Process/Thread
- 1 connection per Thread
- Many threads per Child
- Pros:
 - Efficient use of memory
 - Highly Scalable
- Cons:
 - Faults destroy all threads in that Child

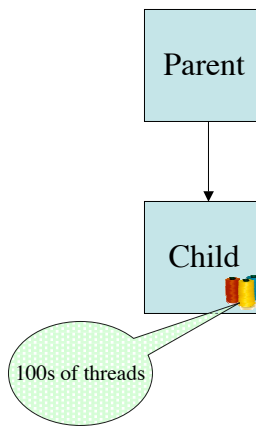


The worker MPM is the new hybrid multitprocess/multithreaded model specifically targeted by Apache 2.0. This MPM trades a small amount of fault isolation for massive scalability (especially in terms of memory consumption). It performs on par or slightly under par with a lightly loaded prefork MPM server, but is able to handle many more concurrent connections without significant memory consumption.

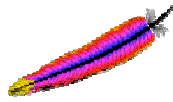


WinNT MPM

- Single Parent/Single Child
- 1 connection per Thread
- Many threads per Child
- Pros:
 - Efficient use of memory
 - Highly Scalable
- Cons:
 - Faults destroy all threads



The WinNT MPM is specifically customized for the Windows platform. It has one parent and one child, with the single child having many hundreds of concurrent threads. Each thread handles a separate request connection.



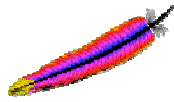
The MPM Breakdown

MPM	Multi-process	Multithreaded
Prefork	Yes	No
Worker	Yes	Yes
WinNT	No*	Yes

* The WinNT MPM has a single parent and a single child.

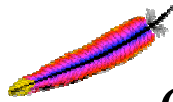
There is some talk of adding multi-child support for the WinNT MPM. This would help isolate faults in 3rd party modules, so that a single failure would not cause new requests to have to wait for a new child to be started (which can mean downtime on the order of seconds, rather than not at all).

Also note that this list is by no means exhaustive. There are a few experimental MPMs that attempt improvements over the worker MPM, as well as a perchild MPM which is targeted at mass multi-hosting facilities.



Efficiency

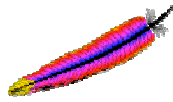
Tricks of the Trade



Questions to ask yourself:

- Do we really need a mutex?
- Can we design so that we don't need it?
- How often will this mutex be locked?
- What is the largest segment size we'll need?

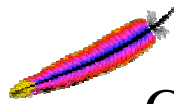
Often times a clever design can avoid requiring a mutex altogether. One good example of this is the scoreboard in Apache. Although the scoreboard is read and written by many processes at the same time, the program was written in such a way to prevent simultaneous use of the same section of the scoreboard.



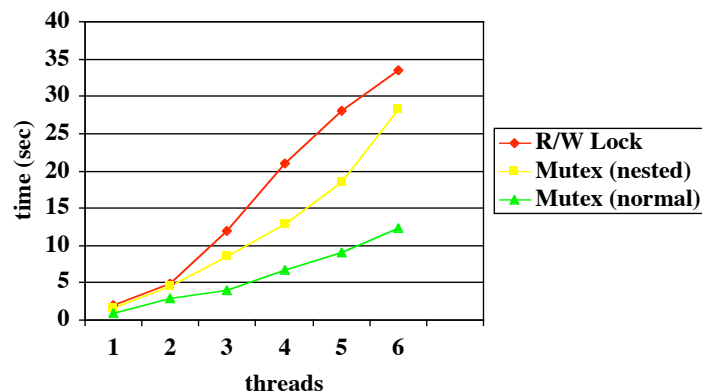
R/W Locks vs. Mutexes

- Reader/Writer locks allow parallel reads
- Mutexes use less overhead
- Reader/Writer locks tend to scale better

Behind the scenes, many rwlock implementations are actually using a mutex and a condition variable or two and a queue of readers and writers in order to resolve contention between the different locking classes. Manipulating those lists can be quite expensive when compared to a simple mutex. Unfortunately APR can not guarantee the same type of performance response across all supported platforms, but it does its best to get out of the way of the OS.



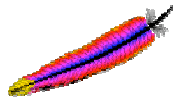
Comparison of Lock Types



In this simple test (which is just a graph of the output from a modified version of the testlockperf program in APR's test directory), we are comparing the nested and non-nested mutexes and read/write locks. At each stage, the test creates a number of threads, and each thread contends for a mutex (the rwlock just grabs the exclusive write lock) increments a counter, and then releases the lock. Once the counter reaches 1 million, the threads are killed and the time is measured.

As is evident in the graph, a rwlock is far more expensive than either mutex type, at least while there are fewer threads contending for the mutex. Also, the nested mutex appears to be on average twice as expensive as the non-nested mutex.

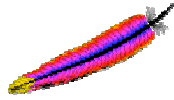
Keep in mind that this is not a very accurate study. I just took my testlockperf and modified it a little so it would run each set of threads for a full test. I will probably commit the changes to CVS. Also keep in mind that I ran this on my uniprocessor laptop, and it is very likely that the results would be very different on other hardware and on another operating system.



Atomics

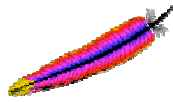
- Pros:
 - Can be very efficient
(sometimes it becomes just one instruction)
- Cons:
 - Produces non-portable binaries
(eg. a solaris 7 binary may not work on solaris 8)

Recently an atomic-lock interface was added to APR. With it came some benefits and drawbacks. Although atomics can provide a very efficient locking mechanism, implementations require intimate knowledge of the underlying system, and often result in non-portable executables.



Threads

Adding threads to your Apache
modules

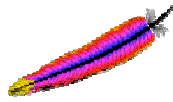


Why use threads?

- background work
- better asynchronous event handling
- pseudo-event-driven models

There are some cases where adding a thread to your module may prove beneficial. Even outside of the realm of Apache modules, you may wish to write a standalone application that takes advantage of APR's thread creation code.

Often times by sticking a thread on an IO-bound operation, one can effectively simulate a fully event-driven model without having to explicitly deal with an event loop. Many thread libraries are very efficient at managing multiple IO-bound threads.



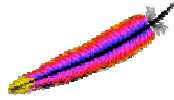
Pitfalls

- pool association
- cleanup registration
- proper shutdown
- async signal handling
- signal masks
- userspace vs. lwp-based thread libraries

Since threads are always associated with a pool, it can be troublesome to synchronize the termination of that pool with the termination of the thread in that pool. APR does **not** terminate threads when their containing pool is cleared or destroyed. This has caused many elusive problems, and is an unfortunate side-effect of tying APR threads to pools.

Often, to maintain a small amount of synchronicity between threads and different events like the shutdown event, it is necessary to register a cleanup routine with the containing pool that somehow signals to the thread that it should shut down. APR does not provide for a way to terminate threads asynchronously, other than terminating the process itself.

There are some platforms that tend to have buggy or incomplete thread libraries. In general, there are very few platforms that behave in the same way, in terms of signal handling, signal masks, and preemption. APR is only able to provide for the lowest-common denominator of abilities, so that it can achieve a common set of portable routines.



Thank You

The End